### 10707 Deep Learning: Spring 2021 Andrej Risteski

Machine Learning Department

Lecture 5: Intro to optimization

### Supervised learning

**Empirical risk minimization approach**: minimize a **training** loss *l* over a class of **predictors**  $\mathcal{F}$ :

$$\hat{f} = \underset{f \in \mathcal{F}}{\operatorname{argmin}} \sum_{\substack{(x,y): \text{training samples}}} l(f(x), y)$$

#### **Three pillars:**

(1) How expressive is the class *F*? (Representational power)

(2) How do we minimize the training loss efficiently? (**Optimization**)

(3) How does  $\hat{f}$  perform on unseen samples? (Generalization)

## The world of continuous optimization

The typical training task in ML can be cast as:  $\min_{x \in \mathbb{R}^d} f(x)$ 

Usually, it is cheap to calculate f(x),  $\nabla f(x)$ , but (more) expensive to calculate higher-order derivatives.

Most algorithms we will look at are iterative: they progressively pick points  $x_1, x_2, ...$  that are supposed to bring "improvement".

Non-exhaustive coverage: entire field of optimization, with applications vastly beyond ML. We focus on deep-learning-relevant methods.

# The mother of all optimization algorithms: gradient descent

The simplest optimization algorithm: Taylor expand and find the direction of "steepest" descent. More precisely:

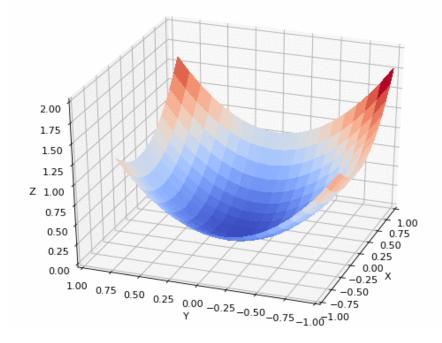
By Taylor's theorem, we have  $f(x + \Delta) \approx f(x) + \Delta^T \nabla f(x) + O(||\Delta||^2)$ 

So, if we ignore higher-order effects, we have

$$\underset{\Delta, ||\Delta|| \le \epsilon}{\operatorname{argmin}} \left\{ f(x + \Delta) - f(x) \right\} = -\epsilon \frac{\nabla f(x)}{||\nabla f(x)||}$$

i.e. we should move (appropriately scaled) opposite of the gradient

### Gradient descent, pictorially

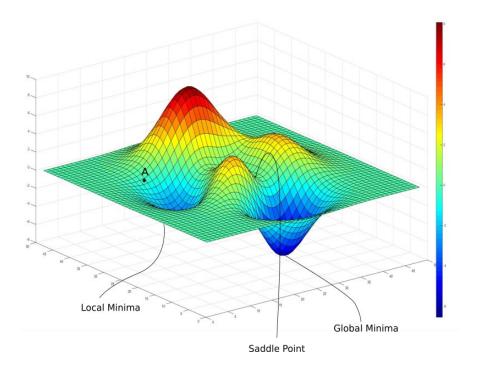


What can we hope for, in the case that  $||\Delta|| \rightarrow 0$ ?

We stop moving when  $\nabla f(\hat{x}) \approx 0$ : these care called **stationary points**.

What kinds of stationary points are there?

## Types of stationary points



**Global minimum**: actual minimizer, namely  $f(\hat{x}) \leq f(x), \forall x \in \mathbb{R}^d$  **Local minimum**:  $f(\hat{x}) \leq f(x), \forall x \text{ s. t. } ||x - \hat{x}|| \leq \epsilon$  for some  $\epsilon > 0$  **Local maximum**:  $f(\hat{x}) \geq f(x), \forall x \text{ s. t. } ||x - \hat{x}|| \leq \epsilon$  for some  $\epsilon > 0$ **Saddle points**: stationary point that is \*not\* a local min/max.

## Types of stationary points Local Minima Global Minima Saddle Point

**Global minimum**: finding these in general is very hard (both in theory – NP-hard, as well as in practice)

**Local minimum**: seem to work quite well often. Some theoretical understanding of why in very restricted cases.

**Saddle points**: typically bad, arise from invariances in input. Want to avoid these. (Stay tuned.)

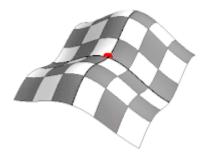
## Checking for local minima?

**Second order checks**: Hessian approximates a function to second order

**Taylor's thm**:  $f(x + \Delta) \approx f(x) + \Delta^T \nabla f(x) + \frac{1}{2} \Delta^T \nabla^2 f(x) \Delta + O(||\Delta||^3)$  $\approx f(x) + \frac{1}{2} \Delta^T \nabla^2 f(x) \Delta + O(||\Delta||^3)$ 



If  $\nabla^2 f(x) > 0$ : for any direction  $\Delta$ , and small enough  $||\Delta||$  $\Delta^T \nabla^2 f(x) \Delta + O(||\Delta||^3) \ge 0$ , so  $f(x + \Delta) > f(x)$ **Local minimum**! (Flipped for local maximum)



If  $\nabla^2 f(x)$  has both positive and negative eigenvalues: Saddle point (not a local minimum/maximum)

If neither of these attains, test is inconclusive!

## The descent lemma: analyzing gradient descent

So far, we've only considered the limit  $||\Delta|| \rightarrow 0$ .

If  $||\Delta||$  is too large, the Taylor expansion will be invalid (and gradient descent can "jump over" local minima).

If  $||\Delta||$  is too small, the runtime of the algorithm will suffer. The descent lemma characterizes the "sweet spot":

**Theorem (descent lemma)**: Let f be twice differentiable, and  $||\nabla^2 f(x)||_2 \le \beta$ . Then, setting  $\eta = 1/\beta$ , and calling  $x_t$  the iterates of gradient descent, namely  $x_{t+1} = x_t - \eta \nabla f(x_t)$ , we have:

$$f(x_t) - f(x_{t+1}) \ge \frac{1}{2\beta} ||\nabla f(x_t)||_2^2$$

### Using the descent lemma: Lyapunov functions

**Theorem (descent lemma)**: Let f be twice differentiable, and  $||\nabla^2 f(x)||_2 \le \beta$ . Then, setting  $\eta = 1/\beta$ , and calling  $x_t$  the iterates of gradient descent, namely  $x_{t+1} = x_t - \eta \nabla f(x_t)$ , we have:

$$f(x_t) - f(x_{t+1}) \ge \frac{1}{2\beta} ||\nabla f(x_t)||_2^2$$

Suppose *f* is lower bounded (e.g.  $f \ge 0$ ), and  $f(x_0) \le M$ Suppose we want point  $x_t$ , s.t.  $||\nabla f(x_t)|| \le \epsilon$ .

**Lyapunov (potential) fn argument**: suppose  $\forall t \in [0, T], ||\nabla f(x_t)|| \ge \epsilon$ Then,  $f(x_T) \le f(x_0) - T \frac{1}{2\beta} \epsilon^2 \le M - T \frac{1}{2\beta} \epsilon^2$ . Also,  $f(x_T) \ge 0$ . Putting these together, we get  $T \le 2M\beta/\epsilon^2$ 

### Proving the Descent Lemma

**Theorem (descent lemma)**: Let *f* be twice differentiable, and  $||\nabla^2 f(x)||_2 \leq \beta$ . Then, setting  $\eta = 1/\beta$ , and calling  $x_t$  the iterates of gradient descent, we have:

$$f(x_t) - f(x_{t+1}) \ge \frac{1}{2\beta} ||\nabla f(x_t)||_2^2$$

**Proof:** By Taylor expansion and the mean value theorem, we have

$$f(x + \Delta) = f(x) + \Delta^T \nabla f(x) + \frac{1}{2} \Delta^T \nabla^2 f(y) \Delta$$

Moreover,  $\Delta^T \nabla^2 f(y) \Delta \leq ||\nabla^2 f(y)||_2 ||\Delta||_2^2 \leq \beta ||\Delta||_2^2$ . Plugging in  $\Delta = -\eta \nabla f(x_t)$ :

$$f(x_{t+1}) \le f(x_t) - \eta \left\| |\nabla f(x_t)| \right\|^2 + \frac{1}{2} \beta \eta^2 \left\| |\nabla f(x_t)| \right\|^2$$
  
=  $f(x_t) - 1/\beta \left\| |\nabla f(x_t)| \right\|^2 + \frac{1}{2} 1/\beta \left\| |\nabla f(x_t)| \right\|^2$   
=  $f(x_t) - 1/2\beta \left\| |\nabla f(x_t)| \right\|^2$ 

### Choosing a learning rate in practice

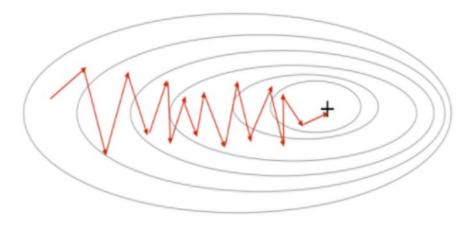
The previous lemma was an instance of a **fixed** learning rate  $(1/\beta)$ . In fact, proof works up to "critical" value of  $2/\beta$ .

**Fixed** learning rates are common: typically grid search b/w  $10^{-6}$  and 1. Pick largest rate for which training doesn't diverge. Frequently ~  $10^{-2}$  works ok.

**Decaying** learning schedules are also popular, e.g. :  $\eta_t = \frac{\eta_0 \tau}{\max(t, \tau)}$ 

(i.e. set learning rate to  $\eta_0$  for the first  $\tau$  iteration, then decay like 1/t) Some intuitions for a decaying rate come from standard convex optimization.

Let's consider *f*'s that are quadratic. (Close to local minima, this will be "true" due to Taylor). What quadratics are bad/good for gradient descent?



**Bad behavior**: gradients don't point towards minimizer – a lot of zigzaging until we reach minimizer.

**Intuitively**: ellipsoidal contours (level sets) should be worse than spherical level sets.

**Question**: Let  $f(x) = \frac{1}{2}x^T A x$ , can we characterize convergence time of gradient descent more precisely? What does it depend on?

**Question**: Let  $f(x) = \frac{1}{2}x^T A x$ , can we characterize convergence time of gradient descent more precisely? What does it depend on?

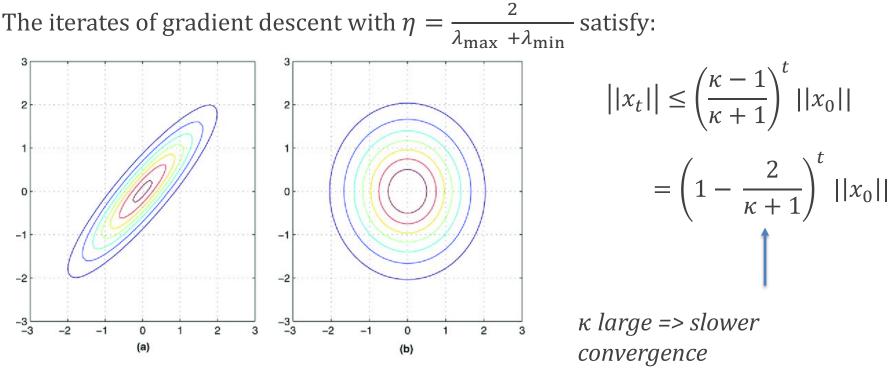
**Thm**: Let *A* be a symmetric positive-definite matrix with minimum and maximum eigenvalues  $\lambda_{\min}$  and  $\lambda_{\max}$  and denote  $\kappa = \lambda_{\max} / \lambda_{\min}$  (condition number).

The iterates of gradient descent with 
$$\eta = \frac{2}{\lambda_{\max} + \lambda_{\min}}$$
 satisfy:  
 $||x_t|| \le \left(\frac{\kappa - 1}{\kappa + 1}\right)^t ||x_0||$ 

 $= ||x_t - 0||$ , *i.e. distance from optimum*  $= ||x_0 - 0||$ , *i.e. distance from optimum* 

**Question**: Let  $f(x) = \frac{1}{2}x^T A x$ , can we characterize convergence time of gradient descent more precisely? What does it depend on?

**Thm**: Let *A* be a symmetric positive-definite matrix with minimum and maximum eigenvalues  $\lambda_{\min}$  and  $\lambda_{\max}$  and denote  $\kappa = \lambda_{\max} / \lambda_{\min}$  (condition number).



**Question**: Let  $f(x) = \frac{1}{2}x^T A x$ , can we characterize convergence time of gradient descent more precisely? What does it depend on?

**Thm**: Let *A* be a symmetric positive-definite matrix with minimum and maximum eigenvalues  $\lambda_{\min}$  and  $\lambda_{\max}$  and denote  $\kappa = \lambda_{\max} / \lambda_{\min}$  (condition number).

The iterates of gradient descent with  $\eta = \frac{2}{\lambda_{\max} + \lambda_{\min}}$  satisfy:  $||x_t|| \le \left(\frac{\kappa - 1}{\kappa + 1}\right)^t ||x_0||$ 

Proof: 
$$||x_{t+1}|| = ||x_t - \eta \nabla f(x_t)||$$
  
 $= ||x_t - \eta A x_t|| = ||(I - \eta A) x_t|| \le ||I - \eta A||_2 ||x_t||_2$   
 $\le \max(|1 - \eta \lambda_{max}|, |1 - \eta \lambda_{min}|) ||x_t||_2$   
 $= \frac{\lambda_{max} - \lambda_{min}}{\lambda_{max} + \lambda_{min}} ||x_t||_2 = \frac{\kappa - 1}{\kappa + 1} ||x_t||_2$ 

## Fixes to the conditioning problem

What can we do for poorly conditioned problems?

Quadratic problem suggests solution: we can solve it in closed form!!

If  $f(x) = \frac{1}{2}x^T A x + b^T x + c$ , minimizer is  $A^{-1}b$ . (Just take derivatives, set to 0.)

What do we do for arbitrary f? Approximate function to second order!!

By Taylor's thm:  $f(x + \Delta) \approx f(x) + \Delta^T \nabla f(x) + \frac{1}{2} \Delta^T \nabla^2 f(x) \Delta + O(||\Delta||^3)$ 

Ignoring 3<sup>rd</sup> and higher order terms, and using the above observation for quadratics:

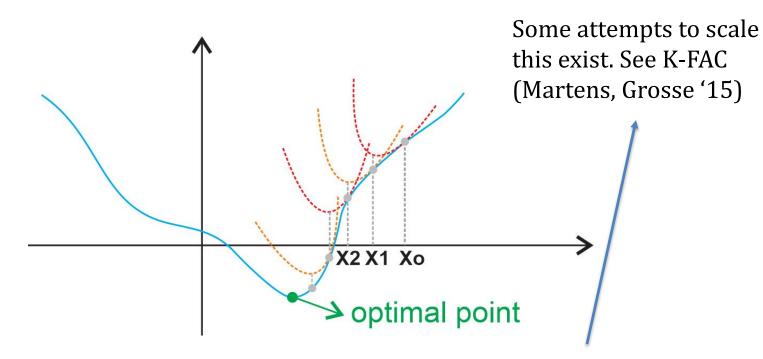
Set 
$$x_{t+1} = x_t - \eta \left( \nabla^2 f(x_t) \right)^{-1} \nabla f(x_t)$$

Newton's method.

### Fixes to the conditioning problem

Set 
$$x_{t+1} = x_t - \eta \left( \nabla^2 f(x_t) \right)^{-1} \nabla f(x_t)$$

Newton's method.



**Problem**: need to invert a  $d \times d$  matrix =>  $d^3$ runtime. Way too expensive.

## Momentum (Polyak '64)

**Alternative fix**: instead of using the gradient at the current step, use a linear combination of the gradients at prior steps. "Smooths" out zig-zagging, by not relying too much on current gradient.

*Linear combination of prior gradients + current one* 

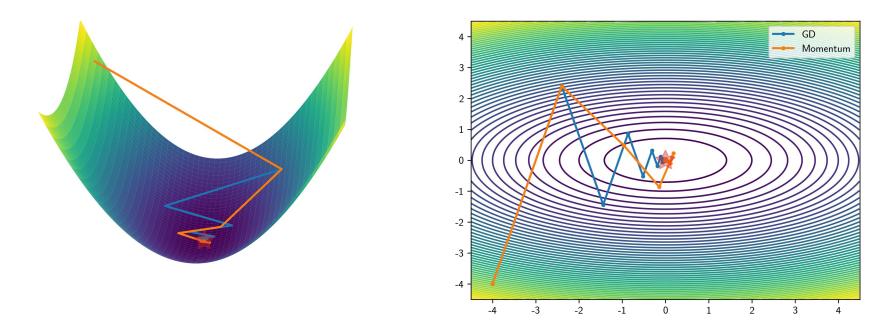
$$v_{t+1} = -\nabla f(x_t) + \beta v_t$$
$$x_{t+1} = x_t + \eta v_{t+1}$$

**Helps provably!** For quadratic case, i.e.  $f(x) = x^T A x$ , you can show:

$$\left||x_t|\right| \le \left(\frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1}\right)^t \left||x_0|\right|$$

It's common in practice to take  $\beta = 0.5, 0.9, 0.99.$ 

## Momentum (Polyak '64)



#### Momentum vs grad. descent https://tangbinh.github.io/01/04/Optimizers.html

## Momentum (Nesterov '83)

**Nesterov acceleration** is a *lookahead* variant of momentum, which has provable benefits for \*any\* convex function. (And is in a certain precise sense, the optimal first-order optimization algorithm).

Evaluate gradient at a  
"lookahead" point  

$$v_{t+1} = -\nabla f(x_t + \beta v_t) + \beta v_t$$
  
 $x_{t+1} = x_t + \eta v_{t+1}$ 

**Magical!** There's been a mini cottage industry to "explain" Nesterov acceleration.

The workhorse for training neural networks: an algorithm that for a network with V nodes and E edges calculates the gradient in **linear time** O(V+E).

The name **backpropagation** was introduced by *Rumelhart, Hinton, Williams '86*, but so natural that it was rediscovered multiple times (as early as 60s). Algorithm seems to first be mentioned in *Werbos'* thesis '74 in the context of neural networks.

In **control theory**: Kelley '60, Bryson '61 [cast as dynamic programming];

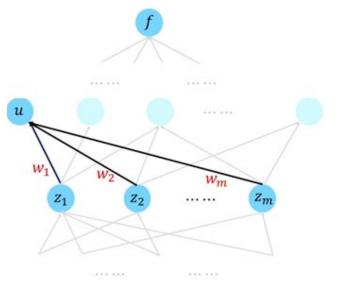
In **theoretical computer science**: Baur-Strassen lemma '83 [in the context of algebraic circuits]

The main tool for deriving backprop: chain rule

Suppose 
$$f(y) = f(x_1(y), x_2(y), ..., x_n(y))$$
  
Then,  $\frac{\partial f}{\partial y} = \sum_{i=1}^{n} \frac{\partial f}{\partial x_i} \frac{\partial x}{\partial y}$ 

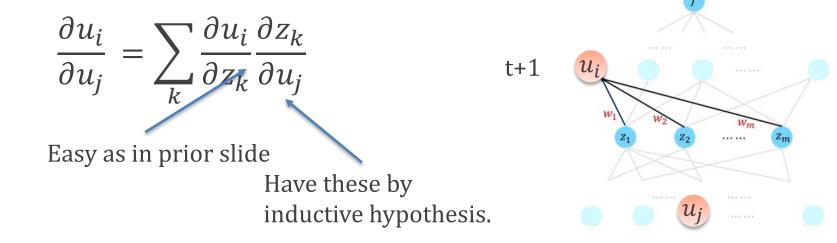
**Observation 1**: It suffices to take derivatives with respect to node functions.

$$\frac{\partial f}{\partial w_1} = \frac{\partial f}{\partial u} \frac{\partial u}{\partial w_1} = \frac{\partial f}{\partial u} \frac{\partial \sigma(\langle w, z \rangle + b)}{\partial w_1}$$
$$= \frac{\partial f}{\partial u} \sigma'(u) z_1$$



**Observation 2**: The obvious forward propagation algorithms results in runtime of  $\Omega(V^2)$ . (Bad! We want O(V+E))

Obvious algorithm? Calculate inductively  $\frac{\partial u_i}{\partial u_j}$ , for all pairs  $(u_i, u_j)$  where  $u_j$  is lower than  $u_i$  (obviously, this includes  $\frac{\partial f}{\partial u}$  which we want)

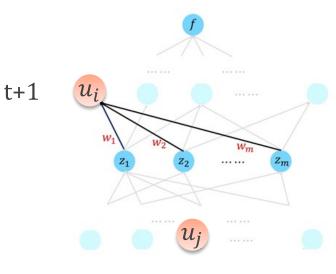


**Observation 2**: The obvious forward propagation algorithms results in runtime of  $\Omega(V^2)$ . (Bad! We want O(V+E))

Obvious algorithm? Calculate inductively  $\frac{\partial u_i}{\partial u_j}$ , for all pairs  $(u_i, u_j)$  where  $u_j$  is lower than  $u_i$  (obviously, this includes  $\frac{\partial f}{\partial u}$  which we want)

$$\frac{\partial u_i}{\partial u_j} = \sum_k \frac{\partial u_i}{\partial z_k} \frac{\partial z_k}{\partial u_j}$$

**Bad** – this will end up with  $\Omega(V^2)$  algorithm.



**Observation 3**: The better way to do this is in a backward fashion.

**Message passing algorithm [dynamic programming]:** each node *u* receives messages (real numbers) from its neighbors on top. Let their sum be *S*. The node passes to downward neighbors z:  $S \frac{\partial u}{\partial z}$ . Proceed from top to down.

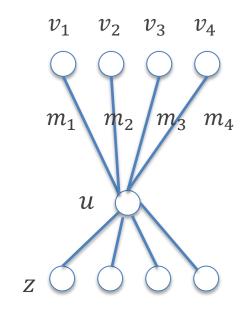
**Claim:** The sum S of the messages that each node u computes is equal to  $\frac{\partial f}{\partial u}$ .

**Proof:** By induction.

Suppose u is at layer t, and inductive hypothesis holds for layers t+1 and above. Sum of messages to u satisfies:

$$S = \sum_{k} m_{k} = \sum_{k} \frac{\partial f}{\partial v_{i}} \frac{\partial v_{i}}{\partial u} = \frac{\partial f}{\partial u}$$

Inductive hypothesis Def. of messages



**Observation 3**: The better way to do this is in a backward fashion.

**Message passing algorithm [dynamic programming]:** each node *u* receives messages (real numbers) from its neighbors on top. Let their sum be *S*. The node passes to downward neighbors z:  $S \frac{\partial u}{\partial z}$ . Proceed from top to down.

**Amount of work:** each node u needs to sum its upward neighbor messages (at most deg(u) of them), and pass a message to its downward neighbors (at most deg(u) of them).

Each downward message just takes an extra  $\frac{\partial u}{\partial z}$  calculation (easy const. time), so each node does O(deg(u)) amount of work.

Hence, **total amount of work** for all nodes is  $O(\sum_u \deg(u)) = O(E)$ 

**Amount of memory:** for calculating  $\frac{\partial u}{\partial z}$ , we need the activation values of intermediate nodes – so **memory** O(V). [Important! If recalculating these, runtime would be quadratic]

